

- Theoretical: max of {backward "distance" / previous ref}
- *Practical*: **min** of {"*time*" to the *most recent* reference}

LRU Timestamp

Choice of implementation

- Absolute timestamp: CPU clock or counter
 - Require additional memory write to the page table (by hardware) on **every memory reference**
- Relative timestamp: linked list of page numbers (hard to do in hardware)
 - Associated with each process is a list of page numbers
 - When a page is referenced, move the corresponding page number to the **head** end
 - \circ $\;$ Victime page is the page at the tail end of the list

LRU Approximations

- Pure implementation of LRU requires accurate timestamp
 - Updating an N-bit timestamp is costlier than updating a 1-bit data
- Approximation Techniques (approximate timestamp with tags)
 - $\circ~$ Bit in tags are set to one by HW, reset to zero by SW
 - k-bit history tag
 - Second-chance algorithm
 - Enhanced second-chance algorithm

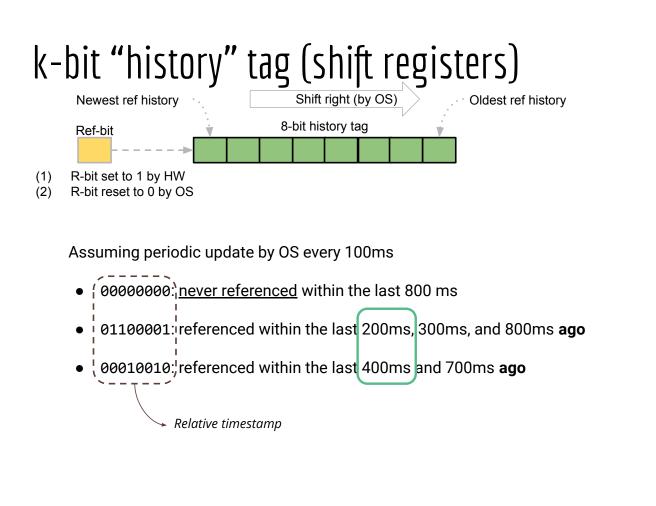
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Approximation #1: reference-bit

- Use a reference-bit in place of actual (absolute) timestamp
 - Include the ref-bit on each cell/row of page table (in addition to the *modified-bit*)
- Updating the ref-bit
 - $\circ~$ Set to one by $\ensuremath{\textbf{hardware}}$ on every memory reference
 - \circ $\;$ Reset to zero by ${\color{black} \textbf{software}}$ (the OS) when the page is swapped in
- Victim Selection: page whose ref-bit is zero
 - Cons: can't distinguish reference 5ms (larger timestamp) vs 20ms (smaller timestamp) ago

Approximation #2: *k-bit "timestamp"*

- Use a **k-bit history tag** (in addition to the ref-bit) for better approximation of timestamp
 - Hardware: the 1-bit ref is automatically set to ONE by MCU
 - **Software**: the OS *periodically*:
 - shifts right the k-bit history and inject the ref-bit at the MSB position of the k-bit history
 - Reset the ref-bit
- Victim Selection: page whose k-bit history is lowest (treated as unsigned integer)
 - Use FIFO to break tie



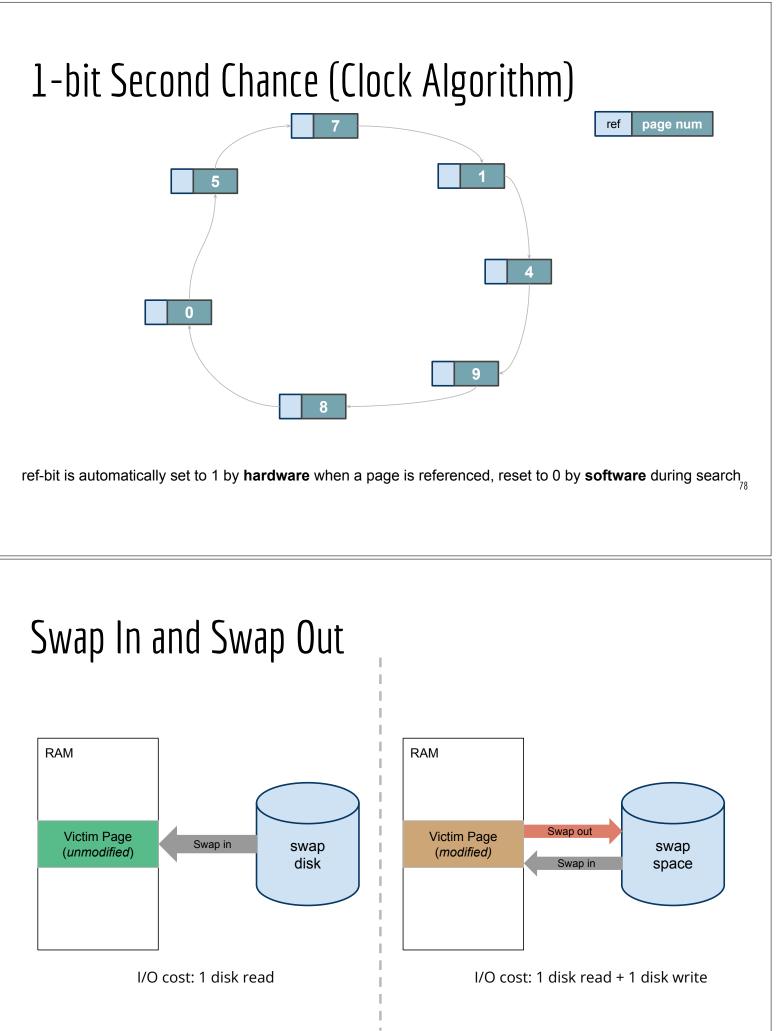
Approximation #3:1-bit Clock Algorithm

• Based on FIFO (uses a circular buffer to maintain the resident page numbers)

Algorithm

- Hardware: set Ref-bit to ONE on each memory reference
- Software
 - Advance the buffer pointer to locate a page whose ref-bit is ZERO
 - While advancing the pointer, reset non-zero reference-bit to ZERO
 - The next search begins at the next position where the last victim was found
- The worst case: all the ref-bits were ONE, the algorithm would have inspected all the pages before finding one victim, giving all the pages a second chance (to stay resident)

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Approximation #4: 2-bit Clock Algorithm

- Improvement to 1-bit clock algorithm, but uses the combination of both ref-bit and modified-bit
- Four types of page

Class	ref-bit	mod-bit	Description
0	0	0	Not (recently)-used and not modified
1	0	1	Not (recently)-used but modified
2	1	0	(recently)-used and not modified
3	1	1	(recently)-used and modified

Victim selection: select from the lowest class available

R-bit

1-bit clock

- Set to 1 by MCU (hardware)
- Reset to 0 by PFI handler (software)
 - At swap in
 - During search

R-bit

2-bit clock

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- Set to 1 by MCU (hardware)
- Reset to 0 by PFI handler (software)
 - At swap in
 - During search

M-bit

- Set to 0 by MCU on memory read
- Set to 1 by MCU on memory write
- Reset to 0 by PFI handler **only** at swap in
- Untouched by PFI handler during search

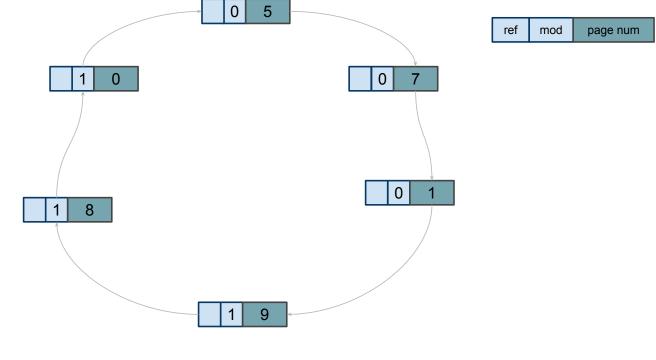
Two-Bit Clock Algorithm: Multi-Pass Search

- Look for Class 0: not-used, not-modified (R:0, M:0)
 - While searching, also reset Ref-bit to zero

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- Class 2 (R=1, M=0) becomes class 0 (R=0, M=0)
- Class 3 (R=1, M=1) becomes class 1 (R=0, M=1)
- If no victim found, search again for "Class 0" (originally were class 2)
- If no victim found (yet), search again for Class 1 (include those originally were class 3)
- Overall effect: modified pages have lower priority to be kicked out

2-bit Second Chance (Enhanced Clock Algorithm)



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